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Fix a real reductive group $G(\mathbb{R})$.

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Atlas of Lie Groups and Representations:

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Atlas of Lie Groups and Representations: Take this idea seriously

p-adic groups

Fix a p-adic group *G*. **Question:** Is there a finite algorithm to compute:

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(So far the answer seems to be no...)

Overview Admissible Dual Unipotent Representations and the Future Three views of the Admissible Dual The Langlands Classification \mathcal{D} -modules K orbits on G/B The Algorithm

Admissible Dual of $G(\mathbb{R})$ Recall $\widehat{G}_u \subset \widehat{G}_a$ (admissible dual)

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```
(2,157 \text{ of them} = .41\% \text{ are unitary})
```

() Input an arbitrary complex reductive algebraic group $G(\mathbb{C})$

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- 1) explicit: a computable combinatorial set
- 2) natural: make the Kazhdan-Lusztig-Vogan polynomials computable

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Three pictures

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Langlands classification: induced from discrete series, characters of Cartan subgroups

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 \mathcal{D} -modules: local systems on $K(\mathbb{C})$ orbits on $G(\mathbb{C})/B(\mathbb{C})$

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L-homomorphism: local systems on the space of admissible homomorphism of the Weil group into the dual group

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For now assume G is simply connected, adjoint and Out(G) = 1(Examples: $G = G_2$, F_4 or E_8)

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The Langlands Classification

Roughly: parametrize representations by characters of Cartan subgroups

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(like the $R_T(\theta)$'s in Deligne-Lusztig's theory for finite groups) Definition:

 $\mathcal{C}(G(\mathbb{R}),\rho) = \{(H(\mathbb{R}),\chi)\}/G(\mathbb{R})$

 $H(\mathbb{R})$ =Cartan subgroup χ = character of $H(\mathbb{R})$ with $d\chi = \rho$

 $(H(\mathbb{R}), \chi) \to I(H(\mathbb{R}), \chi)$ =standard module (induced from discrete series of $M(\mathbb{R})$)

 $\rightarrow \pi(H(\mathbb{R}), \chi)$ (unique irreducible quotient)

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Theorem: The map $(H(\mathbb{R}), \chi) \to \pi(H(\mathbb{R}), \chi)$ induces a canonical bijection:

$$\Pi(G(\mathbb{R}),\rho) \stackrel{1-1}{\longleftrightarrow} \mathcal{C}(G,\rho)$$

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The Langlands Classification

This tells us what we need to compute:

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In particular:

 $|\Pi(G(\mathbb{R}),\rho)| = \sum_{i} |W/W(G(\mathbb{R}),H(\mathbb{R})_{i})||H(\mathbb{R})/H(\mathbb{R})_{i}|$

 $H(\mathbb{R})_1, \ldots, H(\mathbb{R})_n$ are representatives of the conjugacy classes of Cartan subgroups.

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Example: $G(\mathbb{R}) = SL(2, \mathbb{R})$ Overview Admissible Dual Unipotent Representations and the Future K orbits on C For Algorithm

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Example:

 $G(\mathbb{R})=SL(2,\mathbb{R})$

$$A(\mathbb{R}) = \operatorname{diag}(x, \frac{1}{x}) \simeq \mathbb{R}^{\times}, |H(\mathbb{R})/H(\mathbb{R})^{0}| = 2,$$

$$W(G(\mathbb{R}), H(\mathbb{R})) = W = \mathbb{Z}/2\mathbb{Z}$$

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$$T = \begin{pmatrix} \cos(\theta) & \sin(\theta) \\ -\sin(\theta) & \cos\theta \end{pmatrix} \simeq S^1, |H(\mathbb{R})/H(\mathbb{R})^0| = 1, W = 1$$

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 $SL(2, \mathbb{R})$ has 4 irreducible representations of infinitesimal character ρ

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Example: $G = SL(2, \mathbb{R})$, infinitesimal character $=\rho$



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$\mathcal{D} ext{-modules}$

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Definition:

 $\mathcal{D}(G, K, \rho) = \{(x, \chi)\}/K$

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 $x\in \mathcal{B}$

 $\chi =$ local system on $\mathcal{O} = K \cdot x$

= character of $\operatorname{Stab}(x)/\operatorname{Stab}(x)^0$

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Theorem: (Vogan, Beilinson/Bernstein) There is a natural bijection

 $\Pi(G(\mathbb{R}),\rho) \stackrel{1-1}{\longleftrightarrow} \mathcal{D}(G,K,\rho)$

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Example: $G = SL(2, \mathbb{C}), G(\mathbb{R}) = SL(2, \mathbb{R})$ \mathcal{B} is the sphere $= \mathbb{C} \cup \infty$



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$$K \ni z : w \to z^2 w$$

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Three orbits: north pole (0), south pole (∞), open orbit (\mathbb{C}^{\times})

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L-homomorphisms

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L-homomorphisms

Weil group $W_{\mathbb{R}} = \langle \mathbb{C}^{\times}, j \rangle jzj^{-1} = \overline{z}, j^2 = -1$

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Roughly (Langlands): parametrize representations by map of $W_{\mathbb{R}}$ into G^{\vee} (complex dual group)

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Definition:

 $\mathcal{L}(G,\rho) = \{(\phi,\chi)\}/G^{\vee}$

 $\phi: W_{\mathbb{R}} \to G^{\vee}, (\phi(\mathbb{C}^{\times}) \text{ is semisimple, "infinitesimal character } \rho")$ $\chi = \text{local system on } \Omega^{\vee} = G^{\vee} \cdot \phi$ $= \text{character of Stab}(\phi)/\text{Stab}(\phi)^0$
Note: different real forms of *G* all have the same G^{\vee} (no *K* here). This version must take this into account (Vogan's super packets)

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Note: different real forms of *G* all have the same G^{\vee} (no *K* here). This version must take this into account (Vogan's super packets)

Theorem: There is a natural bijection

$$\coprod_{i} \Pi(G(\mathbb{R})_{i},\rho) \stackrel{1-1}{\longleftrightarrow} \mathcal{L}(G,\rho)$$

where $G_1(\mathbb{R}), \ldots, G_n(\mathbb{R})$ are the real forms of *G*. (this version: book by A/Barbasch/Vogan)

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Recapitulation

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Recapitulation

(1) Character Data (orbits of $G(\mathbb{R})$ on Cartans):

$$\Pi(G(\mathbb{R}),\rho) \stackrel{1-1}{\longleftrightarrow} \mathcal{C}(G(\mathbb{R})) = \{(H(\mathbb{R}),\chi)\}/G(\mathbb{R})$$

(2) \mathcal{D} -modules (orbits \mathcal{O} of K on G/B):

$$\Pi(G(\mathbb{R}),\rho) \stackrel{1-1}{\longleftrightarrow} \mathcal{D}(G,K,\rho) = \{(x,\chi)\}/K$$

(3) L-homomorphisms (orbits Ω^{\vee} of G^{\vee} on L-homomorphisms):

$$\prod_{i=1}^{n} \Pi(G(\mathbb{R})_{i},\rho) \stackrel{1-1}{\longleftrightarrow} \mathcal{L}(G,\rho) = \{(\phi,\chi)\}/G^{\vee}$$

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In each case there is some geometric data, and a character of a finite abelian group (two-group)

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We'd rather talk about orbits than characters of $(\mathbb{Z}/2\mathbb{Z})^n$ (Matching the three pictures: easy up to χ)

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Drop the χ 's and get sets of representations:

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Definition: Orbit Ω^{\vee} of G^{\vee} on L-homomorphisms \rightarrow L-packet

$\Pi_L(G(\mathbb{R}), \Omega^{\vee})$

(or $\coprod_i \Pi_L(G(\mathbb{R})_i, \Omega))$

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(or $\coprod_i \Pi_L(G(\mathbb{R})_i, \Omega)$) **Definition**: Orbit \mathcal{O} of K on $G/B \to \mathbb{R}$ -packet

 $\Pi_R(G(\mathbb{R}),\mathcal{O})$

Theorem (Vogan): The intersection of an L-packet and an R-packet is at most one element.

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Corollary: $\Pi(G(\mathbb{R}), \rho)$ is parametrized by a subset of pairs

(*K* orbit on \mathcal{B} , G^{\vee} orbit on L-homomorphisms)

via

 $(\mathcal{O}, \Omega^{\vee}) \to \Pi_R(G(\mathbb{R}), \mathcal{O}) \cap \Pi_L(G(\mathbb{R}), \Omega^{\vee})$

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Which pairs?...

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K-orbits on the dual side

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Proposition: There is a natural bijection:

$$\mathcal{L} \longleftrightarrow^{1-1} \prod_{i=1}^n K_i^{\vee} \backslash \mathcal{B}^{\vee}$$

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Symmetric Picture

Corollary: $\Pi(G(\mathbb{R}), \rho)$ is parametrized by a subset of pairs

(*K* orbit on \mathcal{B}, K^{\vee} orbit on \mathcal{B}^{\vee})

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Note: This symmetry is Vogan Duality.

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Note: This symmetry is Vogan Duality.

This reduces the problem to:

Parametrize *K* orbits on $\mathcal{B} = G/B$

(applied to G and G^{\vee})

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K orbits on G/B

Definition:

$$\mathcal{X} = \{x \in \operatorname{Norm}_G(H) \mid x^2 = 1\}/H$$

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Proposition: There is a natural bijection

$$\mathcal{X} \stackrel{1-1}{\longleftrightarrow} \coprod_i K_i \backslash \mathcal{B}$$

(union over real forms, corresponding K_1, \ldots, K_n)

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Sketch of Proof

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Sketch of Proof

$$\mathcal{P} = \{(x, B)\}/G \ (x^2 = 1, B = \text{Borel})$$



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Fix representatives x_1, \ldots, x_n of \mathcal{X}/G (i.e. real forms) Fix $B_0 \supset H$

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 $(x, B) \sim_G (x', B_0) \to x' \in \mathcal{X} \quad (\text{wlog } x' \in \text{Norm}(H))$

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$K \setminus G/B$ for $Sp(4, \mathbb{R})$ and SO(3, 2):

 $Sp(4, \mathbb{R})$:

0:	1	2	6	4	[nn]	0	0	
1:	0	3	6	5	[nn]	0	0	
2:	2	0	*	4	[cn]	0	0	
3:	3	1	*	5	[cn]	0	0	
4:	8	4	*	*	[Cr]	2	1	2
5:	9	5	*	*	[Cr]	2	1	2
6:	6	7	*	*	[rC]	1	1	1
7:	7	6	10	*	[nC]	1	2	2,1,2
8:	4	9	*	10	[Cn]	2	2	1,2,1
9:	5	8	*	10	[Cn]	2	2	1,2,1
10:	10	10	*	*	[rr]	3	3	1,2,1,2

SO(3, 2):

0:	0	1	3	2	[nn]	0	0	
1:	1	0	*	2	[cn]	0	0	
2:	5	2	*	*	[Cr]	2	1	2
3:	3	4	*	*	[rC]	1	1	1
4:	4	3	б	*	[nC]	1	2	2,1,2
5:	2	5	*	б	[Cn]	2	2	1,2,1
6:	б	б	*	*	[rr]	3	3	1,2,1,2

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Structure of G

 \mathcal{X} gives structure of G: real forms, Cartan subgroups, Weyl groups

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 \mathcal{X} gives structure of G: real forms, Cartan subgroups, Weyl groups (assume G is adjoint, inner class of compact group)

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Structure of G

 \mathcal{X} gives structure of G: real forms, Cartan subgroups, Weyl groups

(assume G is adjoint, inner class of compact group)

Proposition

- 1) Real forms of $G \stackrel{1-1}{\longleftrightarrow} \mathcal{X}_1 / W$ (\mathcal{X}_1 = fiber over $1 \in W$)
- 2) Cartan subgroups in all real forms: \mathcal{X}/W
- 3) $W(G(\mathbb{R}), H(\mathbb{R})) = \operatorname{Stab}_W(x)$

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The Parameter Space \mathcal{Z}

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The Parameter Space \mathcal{Z}

 $\mathcal{X} \in x$
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The Parameter Space \mathcal{Z}

 $\mathcal{X} \in x \to \Theta_x = \operatorname{int}(x)$

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The Parameter Space \mathcal{Z}

$\mathcal{X} \in x \to \Theta_x = \operatorname{int}(x) \to \Theta_{x,H} = \Theta_x|_{\mathfrak{H}}$

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The Parameter Space \mathcal{Z}

 $\mathcal{X} \in x \to \Theta_x = \operatorname{int}(x) \to \Theta_{x,H} = \Theta_x|_{\mathfrak{H}}$ By symmetry define $\mathcal{X}^{\vee}, \mathcal{X}^{\vee} \ni y \to \Theta_{y,H^{\vee}}$

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Definition:

$$\mathcal{Z} = \{(x, y) \mid \in \mathcal{X} \times \mathcal{X}^{\vee} \mid \Theta_{x, H}^{t} = -\Theta_{y, H^{\vee}} \}$$

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$$\mathcal{Z} = \{(x, y) \mid \in \mathcal{X} \times \mathcal{X}^{\vee} \mid \Theta_{x, H}^{t} = -\Theta_{y, H^{\vee}} \}$$

$$\mathcal{Z} \subset \prod_i K_i ackslash \mathcal{B} imes \prod_j K_j^{ee} ackslash \mathcal{B}^{ee}$$

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The Parameter Space \mathcal{Z}

Theorem: There is a natural bijection:

$$\mathcal{Z} \stackrel{1-1}{\longleftrightarrow} \prod_{i=1}^{n} \Pi(G(\mathbb{R})_{i}, \rho)$$

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Recall $\mathcal{Z} = \{(x, y)\}$

 $x \in \mathcal{X} = \{x \in \operatorname{Norm}_G(H) | x^2 = 1\}/H$ $y \in \mathcal{X}^{\vee} =$ same thing on dual side

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(Canonical up to characters of $G_{qs}(\mathbb{R})/G_{qs}(\mathbb{R})^0$, $G_{qs}^{\vee}(\mathbb{R})/G_{qs}^{\vee}(\mathbb{R})^0$)

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General Groups

General Groups

For simplicity we assumed (recall $G = G(\mathbb{C})$):

- G is simply connected
- **2** G is adjoint
- $\bigcirc \text{Out}(G) = 1$

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For simplicity we assumed (recall $G = G(\mathbb{C})$):

- G is simply connected
- G is adjoint
- $\bigcirc \text{Out}(G) = 1$

In general:

- Fix an inner class of real forms
- **2** Need twists $G^{\Gamma} = G \rtimes \Gamma$, $G^{\vee} \rtimes \Gamma$ ($\Gamma = \text{Gal}(\mathbb{C}/\mathbb{R})$)
- Solution Require $x^2 \in Z(G)$ (not $x^2 = 1$)
- Need several infinitesimal characters
- Need strong real forms

The General Algorithm

$$\mathcal{X} = \{x \in \operatorname{Norm}_{G^{\Gamma} \setminus G}(H) \mid x^2 \in Z(G)\}/H$$

 \mathcal{X}^{\vee} similarly, $\mathcal{Z} = \{(x, y) \mid ...\} \subset \mathcal{X} \times \mathcal{X}^{\vee}$ as before.

Theorem: There is a natural bijection

$$\mathcal{Z} \stackrel{1-1}{\longleftrightarrow} \prod_{i \in S} \Pi(G(\mathbb{R})_i, \Lambda)$$

 Λ = certain set of infinitesimal characters *S* is the set of "strong real forms"

Reference: Algorithms for Representation Theory of Real Reductive Groups, preprint (www.liegroups.org), Fokko du Cloux, A

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Cayley Transforms and Cross Actions

Two natural ways of constructing new representations from old (Vogan): Cayley transforms and cross action

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In our picture:

1) W acts by conjugation on \mathcal{X} and \mathcal{Z} : cross action

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$$w \to w' = s_{\alpha} w \in W_2$$

Two natural ways of constructing new representations from old (Vogan): Cayley transforms and cross action

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1) W acts by conjugation on \mathcal{X} and \mathcal{Z} : cross action

2) $w \in W_2, s_\alpha w = w s_\alpha$,

$$w \to w' = s_{\alpha} w \in W_2$$

lifts to

$$x \to x' = \sigma_a x$$

(Multivalued due to choice of σ_a : x' or $\{x'_1, x'_2\}$) This is the Cayley transform

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 $K \setminus G/B$ for SO(5, 5)

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Closeup of SO(5, 5) graph

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Example: SL(2)/PGL(2)

$PGL(2, \mathbb{C})$:

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Example: SL(2)/PGL(2)

 $PGL(2, \mathbb{C})$:

 $\mathcal{X} = \{I, \operatorname{diag}(-1, -1, 1), w\} \rightarrow$

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 $K(\mathbb{C})$ orbits on $G(\mathbb{C})/B(\mathbb{C})$: { \mathbb{C}^{\times} , ∞ }, { \cdot }

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Overview The L: Admissible Dual D-mo Unipotent Representations and the Future K orbi

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 $SL(2, \mathbb{C}): \mathcal{X} = \{\pm I, \pm \operatorname{diag}(i, -i), w\} \rightarrow$

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Example: SL(2)/PGL(2)

О	x	<i>x</i> ²	K	$G_{X}(\mathbb{R})$	λ	rep	\mathcal{O}^{\vee}	у	y^2	K^{\vee}	$G_y^{\vee}(\mathbb{R})$	λ	rep
•	Ι	Ι	G	SU(2,0)	ρ	\mathbb{C}	$\mathbb{C}^{ imes}$	w	Ι	$O(2,\mathbb{C})$	SO(2, 1)	2ρ	PS_+
	-I	Ι	G	SU(0,2)	ρ	C	$\mathbb{C}^{ imes}$	w	Ι	$O(2,\mathbb{C})$	SO(2, 1)	2ρ	PS_
{0}	t	-I	C×	SU(1,1)	ρ	DS+	$\mathbb{C}^{ imes}$	w	Ι	$O(2,\mathbb{C})$	SO(2, 1)	ρ	\mathbb{C}
$\{\infty\}$	-t	-I	C×	SU(1,1)	ρ	DS_	$\mathbb{C}^{ imes}$	w	Ι	$O(2,\mathbb{C})$	SO(2, 1)	ρ	sgn
$\mathbb{C}^{ imes}$	w	-I	\mathbb{C}^{\times}	SU(1, 1)	ρ	\mathbb{C}	$\{\infty\}$	t	Ι	$O(2,\mathbb{C})$	SO(2, 1)	ρ	DS
$\mathbb{C}^{ imes}$	w	Ι	$O(2,\mathbb{C})$	SU(1, 1)	ρ	PS	•	Ι	Ι	G^{\vee}	SO(3)	ρ	\mathbb{C}

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SL(2)/PGL(2) via atlas output

```
main: type
Lie type: Al sc s
main: block
(weak) real forms are:
0: su(2)
1: sl(2,R)
enter your choice: 1
possible (weak) dual real forms are:
0: su(2)
1: sl(2,R)
enter your choice: 1
entering block construction ...
2
done
Name an output file (return for stdout, ? to abandon):
0(0,1): 1 (2,*) [i1] 0
1(1,1): 0 (2,*) [i1] 0
2(2,0): 2 (*,*) [r1] 1 1
```

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Example: $Sp(4, \mathbb{R})$

```
main: type
Lie type: C2 sc s
main: block
(weak) real forms are:
0: sp(2)
1: sp(1,1)
2: sp(4,R)
enter your choice: 2
possible (weak) dual real forms are:
0: so(5)
1: so(4,1)
2: so(2,3)
enter your choice: 2
entering block construction ...
10
done
Name an output file (return for stdout, ? to abandon):
               2
                   (6, *) (4, *) [i1,i1] 0
0(0,6):
          1
1(1,6):
           0
               3
                   (6, *) (5, *) [i1,i1] 0
2(2,6):
           2 0
                   (*,*) (4,*)
                                      [ic,i1] 0
3(3,6):
           3 1
                     *, *)
                           (5,*)
                                      [ic,i1] 0
4(4,4):
           8
               4
                     *, *) (*, *)
                                      [C+,r1] 1
                                                 2
5(5,4):
           9
               5
                    (*,*) (*,*)
                                      [C+,r1] 1
                                                  2
6(6,5):
           6
               7
                    (*, *) (*, *)
                                      [r1,C+] 1
                                                 1
7(7,2):
           7
              6
                   (10.11)
                           (*, *)
                                      [i2,C-]
                                              2
                                                 2,1,2
8(8,3):
           4
               9
                   (*, *)
                           (10, *)
                                      [C-,i1] 2 1,2,1
               8
                           (10, *)
                                      [C-,i1] 2 1,2,1
9(9,3):
           5
                     *, *)
              10
                                      [r2,r1] 3 1,2,1,2
10(10,0):
          11
                     *, *)
                           (*,*)
11(10.1):
          10
              11
                     *, *)
                            (*, *)
                                      [r2,rn]
                                             3 1.2.1.2
```

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Example: E_8

```
real: type
Lie type: E8 sc s
main: blocksizes
               compact quaternionic split
compact
               0
                        0
                                      1
quaternionic
                        3,150
                                      73,410
               0
split
               1
                        73,410
                                      453,060
```

Proposition: From the output of atlas one can list the special unipotent representations associated to a given nilpotent orbit.

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Sketch

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*) Fix a block ${\mathcal B}$

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Sketch

*) Fix a block \mathcal{B} (block)

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Sketch

*) Fix a block \mathcal{B} (block)

*) Fix nilpotent orbit \mathcal{O} for \mathfrak{g}^{\vee} . Let $S = \{i_1, \ldots, i_r\}$ be the nodes of Dynkin diagram labelled 2. Let $\lambda =$ corresponding infinitesimal character.

Unipotent Representations

1) $\mathcal{O} \to \sigma$ (special representation of *W*)
O → σ (special representation of *W*)
Find all cells *C* ⊂ *B*

- 1) $\mathcal{O} \to \sigma$ (special representation of *W*)
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- 3) List cells *C* containing the special representation $\sigma \otimes \text{sgn}$

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- 4) For each such C list $\pi \in C$ with $\tau(\pi) = S$ (block)

- 1) $\mathcal{O} \to \sigma$ (special representation of *W*)
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(wgraph + calculation with character table of W)

4) For each such C list $\pi \in C$ with $\tau(\pi) = S$ (block)

5) Push these to λ

David Vogan has carried this out for E_8

(70 nilpotent orbits; 20 even ones; 143 unipotent representations with integral infinitesimal character for $E_8(split)$)

Conjecture (Arthur): These representations are unitary.

block dual to split group:			
Rep (x,y)	length	Cartan	roots
133 (133,320205)	0	0	[i1,i1,i1,i1,i1,i1,i1,i1]
140 (140,320204)	1	1	[i1,i1,i1,i1,i1,i1,C+,r1]
42248 (40972,306175)	16	3	[C+,i1,C+,i1,C+,C+,C+,C-]
82083 (77494,287709)	21	6	[C-,C+,rn,C+,rn,C+,rn,C+]
124391(114466,263402)	24	2	[i1,C+,i1,i1,C+,C+,C-,i1]
124432(114507,263398)	24	2	[i1,i1,C+,C+,C-,i1,C+,i1]
132306(120375,257307)	25	3	[C+,C+,i1,i1,C+,C+,C-,C+]
191385(168884,220459)	29	1	[i1,i1,i1,i1,i1,i1,i1,c-]
198367(172894,213960)	30	4	[C-,C+,C+,C+,C+,C+,C+,C+,C+]
205069(179284,210683)	30	2	[r1,i1,C+,i1,i1,i1,i1,C-]
225144(192668,195053)	32	5	[i1,rn,i1,C+,rn,C+,rn,C-]
233376(200324,190190)	32	2	[C-,i1,C+,C-,C+,C+,C-,C+]
233395(200343,190188)	32	2	[C-,C+,i1,i1,i1,C-,C+,C+]
237240(201594,186548)	33	б	[rn,C-,C+,C+,C+,C+,C+,C+,C+]
243756(206740,180794)	33	3	[C+,i1,C+,C+,C-,C+,i1,C+]
244076(207060,180688)	33	3	[C+,C+,C+,C-,C+,C+,C+,C-]
252552(212118,174728)	34	4	[C+,C+,C+,C+,C-,C+,C+,C+]
258013(216823,170023)	34	3	[C+,C+,C+,i1,i1,i1,C-,C+]
258048(216858,170012)	34	3	[C+,i1,i2,C+,C-,C+,i1,C+]
288684(238673,147429)	36	2	[C+,i1,i1,i1,i1,C-,C+,i1]
309166(250360,129909)	38	4	[C+,C-,C+,C+,C+,C+,C+,C+,C-]
320784(257336,120344)	39	4	[C+,C+,C+,i2,C-,C+,C+,C+]
453058(320205, 133)	64	9	[r2,r2,r2,r2,r2,r2,r2,r2]
block dual to compact	group:		
0 (0, 320205)	0	0	[ic,ic,ic,ic,ic,ic,ic,ic]

Overview Admissible Dual Unipotent Representations and the Future

What next?

• Put in λ

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• K-structure of representations

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